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# **EECE 417 Computer Systems Architecture**

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# **Computer Organization and Design (3<sup>rd</sup> Ed)**

**-The Hardware/Software Interface**

**by**

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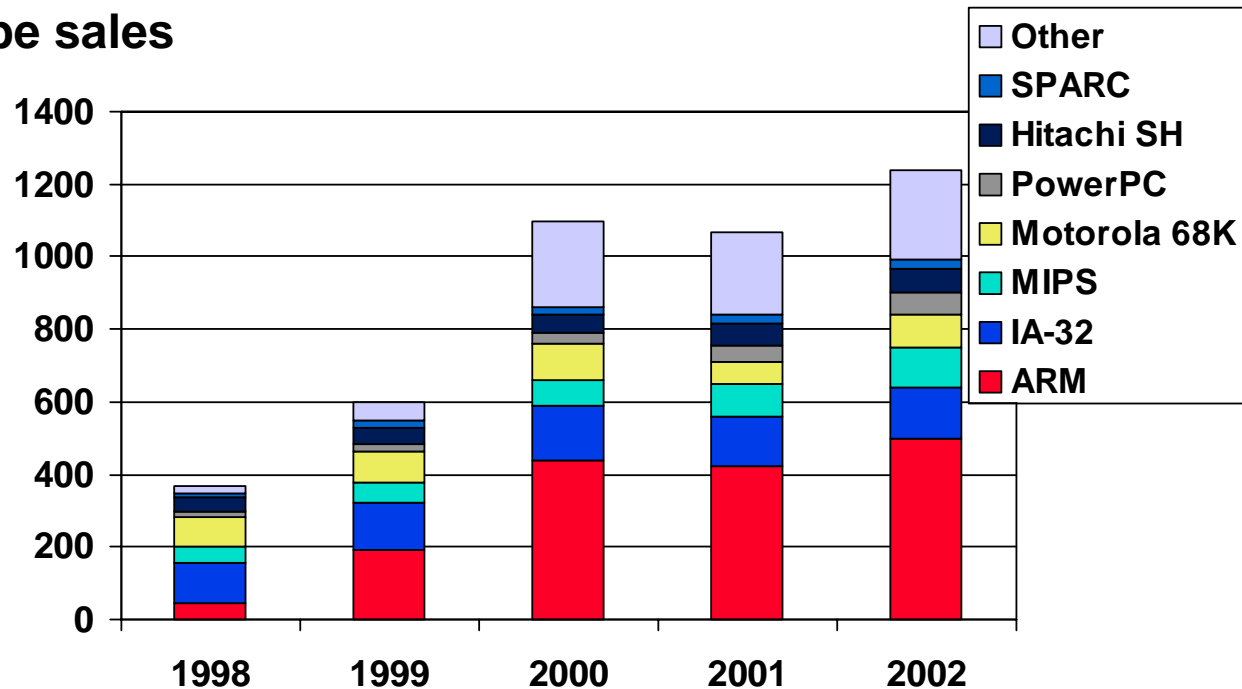
## **Chapter 2**

# **Instructions: Language of the Computer**

# Instructions:

- Language of the Machine
- We'll be working with the MIPS instruction set architecture
  - similar to other architectures developed since the 1980's
  - Almost 100 million MIPS processors manufactured in 2002
  - used by NEC, Nintendo, Cisco, Silicon Graphics, Sony, ...

- ISA type sales



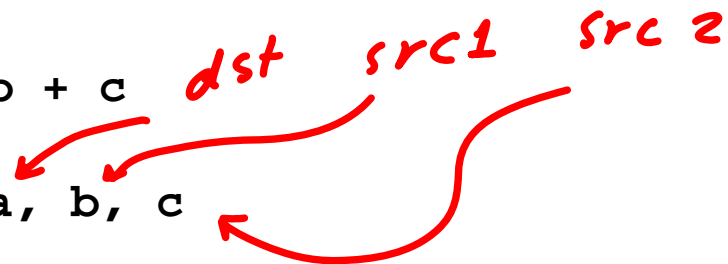
# MIPS arithmetic

- All instructions have 3 operands
- Operand order is fixed (destination first)

Example:

C code:             $a = b + c$

MIPS 'code':      `add a, b, c`



(we'll talk about registers in a bit)

*“The natural number of operands for an operation like addition is **three**...requiring every instruction to have exactly **three operands**, no more and no less, conforms to the philosophy of **keeping the hardware simple**”*

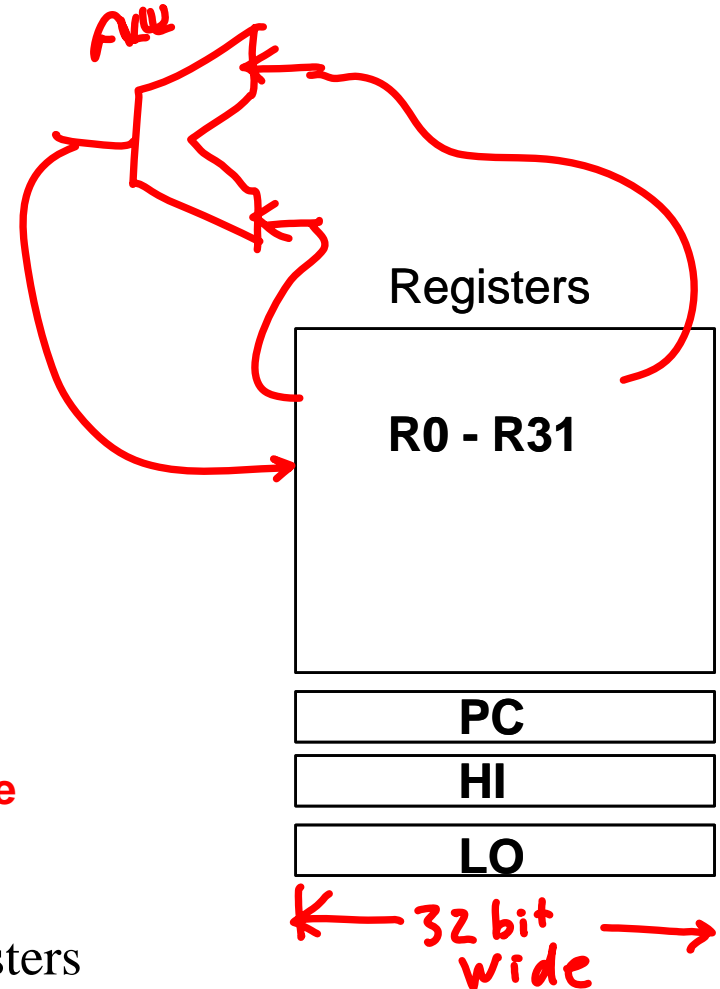
# MIPS arithmetic

- Design Principle: simplicity favors regularity.
- Of course this complicates some things...

C code:            `a = b + c + d;`

MIPS code:        `add a, b, c`  
                     `add a, a, d`

- **Operands must be registers**, only 32 registers provided
- Each register contains 32 bits
- Design Principle: smaller is faster. **Why?**
  - **Large number of registers may increase the clock cycle time**
  - **Then why not 30 registers?**
    - Balance between cries for more registers from programmer vs. fast clock cycle from hardware designer



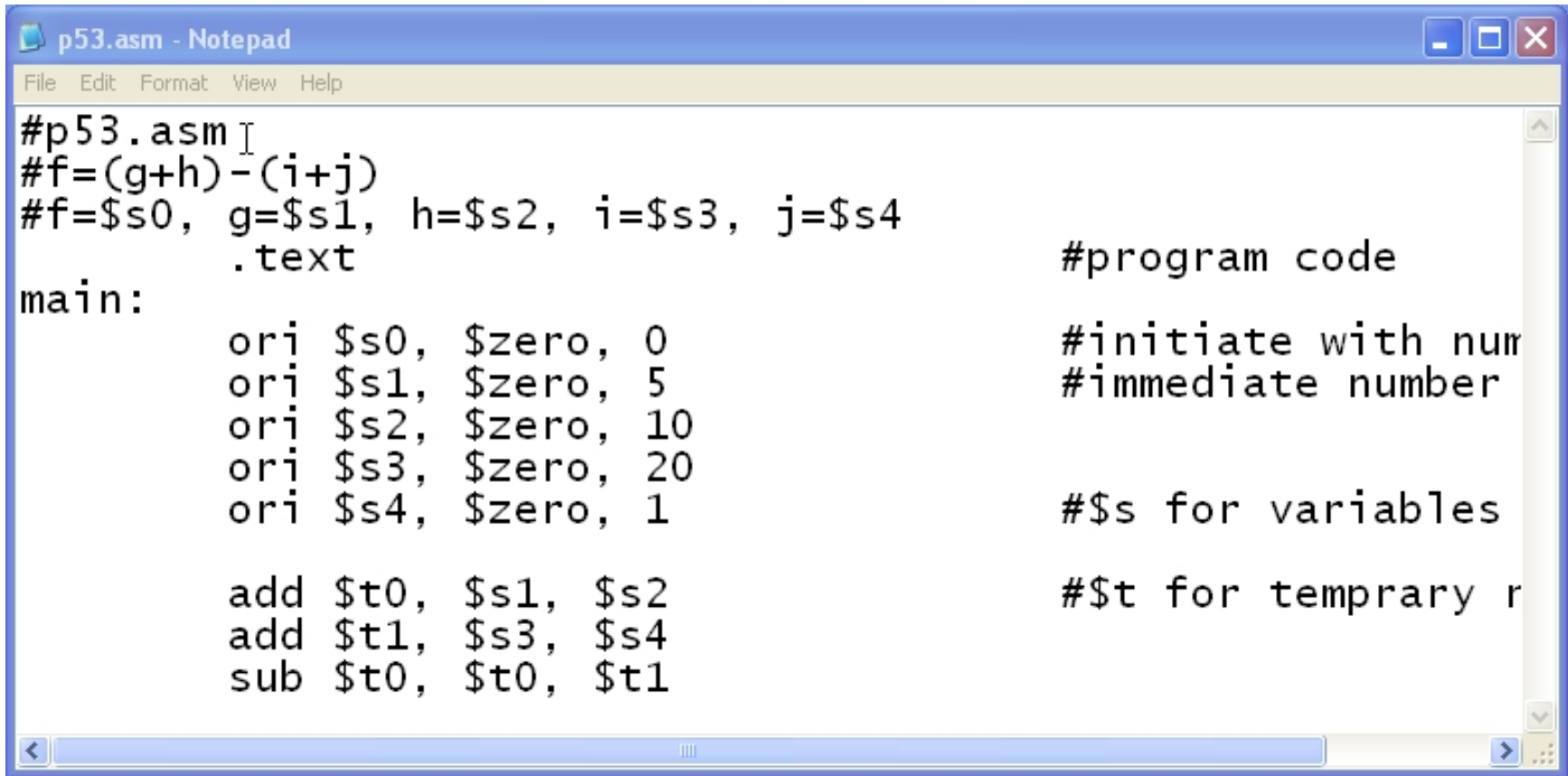
# 32 Registers & Policy of Use Conventions

Register 1 (\$at) reserved for assembler, 26-27 for operating system

Name	Register number	Usage
\$zero	0	the constant value 0
\$at	1	Assembler temporary (reserved)
\$v0-\$v1	2-3	values for results and expression evaluation
\$a0-\$a3	4-7	arguments
\$t0-\$t7	8-15	temporaries
\$s0-\$s7	16-23	saved
\$t8-\$t9	24-25	more temporaries
\$k0-\$k1	26-27	reserved for OS kernel
\$gp	28	global pointer
\$sp	29	stack pointer
\$fp	30	frame pointer
\$ra	31	return address

# Simple Example with MIPS Assembly Language using *SPIM*

- $F = (g+h) - (i+j)$
- $G=5, h=10, i=20, \text{ and } j=1$



```
p53.asm - Notepad
File Edit Format View Help
#p53.asm
#f=(g+h)-(i+j)
#f=$s0, g=$s1, h=$s2, i=$s3, j=$s4
.text
main:
    ori $s0, $zero, 0
    ori $s1, $zero, 5
    ori $s2, $zero, 10
    ori $s3, $zero, 20
    ori $s4, $zero, 1

    add $t0, $s1, $s2
    add $t1, $s3, $s4
    sub $t0, $t0, $t1

#program code
#initiate with nur
#immediate number
#s for variables
#t for temporary r
```

- See the result of **SPIM**. Where is the result?



# SPIM's Human Interface (syscall)

- SPIM provides a small set of operating system-like services through the **system call (syscall)** instruction
- To request a service, a program loads the system call code into register **\$v0** and arguments into registers **\$a0~\$a3**
- System calls that return values put their results in register **\$v0**

Service	System Call Code	Arguments	Result
Print_int	1	\$a0 = integer	
Print_float	2	\$f12 = float	
Print_double	3	\$f12 = double	
Print_string	4	\$a0 = string	
Read_int	5		Integer (in \$v0)
Read_float	6		Float (in \$f0)
Read_double	7		Double (in \$f0)
Read_string	8	\$a0 = buffer, \$a1 = length	
Sbrk	9	\$a0 = amount	Address (in \$v0)
exit	10		

# Simple Example with Syscall

```
p53A.asm - Notepad
File Edit Format View Help
#p53A.asm
#Print out syscall
#f=(g+h)-(i+j)
#f=$s0, g=$s1, h=$s2, i=$s3, j=$s4
        .text                #program code
main:
        ori $s0, $zero, 0    #initiate with number
        ori $s1, $zero, 5    #immediate number loading
        ori $s2, $zero, 10
        ori $s3, $zero, 20
        ori $s4, $zero, 1    #$$s for variables

        add $t0, $s1, $s2    #$$t for temprary results
        add $t1, $s3, $s4
        sub $t0, $t0, $t1

#syscall
#calling system call using $v0
# print intger    $v0=1 with argument in $a0
# read integer    $v0=5 with saved in $v0
        ori $v0, $zero, 1    #request for print
        or  $a0, $zero, $t0   #what is in t0
        syscall
```

- Result

 Console

-6|

# Let's read values from keyboard

```
#p53B.asm
#Read 4 values from keyboard
#And Print Out the result using SYSCALL
#a=b+c
#a=$s0, b=$s1, c=$s2
        .text                                #program code
main:
        ori $s0, $zero, 0                    #initiate with number 0
# Read Input (b)
        ori $v0, $zero, 5
        syscall                               #now type-in is in v0
        or  $s1, $zero,$v0
# Read Input (c)
        ori $v0, $zero, 5
        syscall                               #now type-in is in v0
        or  $s2, $zero,$v0
#Add (b) and (c)
        add $s0, $s1, $s2                    #$t for temprary results
#syscall
#calling system call using $v0
# print intger $v0=1 with argument in $a0
# read integer $v0=5 with saved in $v0
        ori $v0, $zero,1                    #request for print
        or  $a0, $zero, $s0                 #what is in t0
        syscall
```

- **Result**
- **a=b+c**

 **Console**

```
-5
4
-1
```

# Let's add bells and whistles

```
#p53C.asm
#Read 4 values from Keyboard
#with strings
#And Print Out the result using SYSCALL
#a=b+c
#a=$s0, b=$s1, c=$s2
main:
    .data                                #data part
msg1:  .ascii "\nType the first number: " #string to print
msg2:  .ascii "\nType the second number: "
msg3:  .ascii "\nThe answer is: "
    .text                                #code part
ori $v0, $zero, 4                       #msg1
la $a0, msg1                             # takes the address of string as an argument
syscall
ori $v0, $zero, 5                       #read input (b)
syscall                                  #now type-in is in v0
or $s1, $zero, $v0
ori $v0, $zero, 4                       #msg2
la $a0, msg2                             # takes the address of string as an argument
syscall
ori $v0, $zero, 5                       #read input (c)
syscall                                  #now type-in is in v0
or $s2, $zero, $v0
add $s0, $s1, $s2                       #add (b) and (c) $t for temp
ori $v0, $zero, 4                       #msg3
la $a0, msg3                             # takes the address of string as an argument
syscall
ori $v0, $zero, 1                       #request for print
or $a0, $zero, $s0                      #result
syscall
```

Console

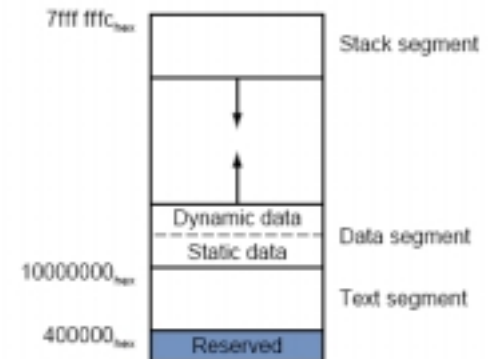
Type the first number: -3

Type the second number: 9

The answer is: 6

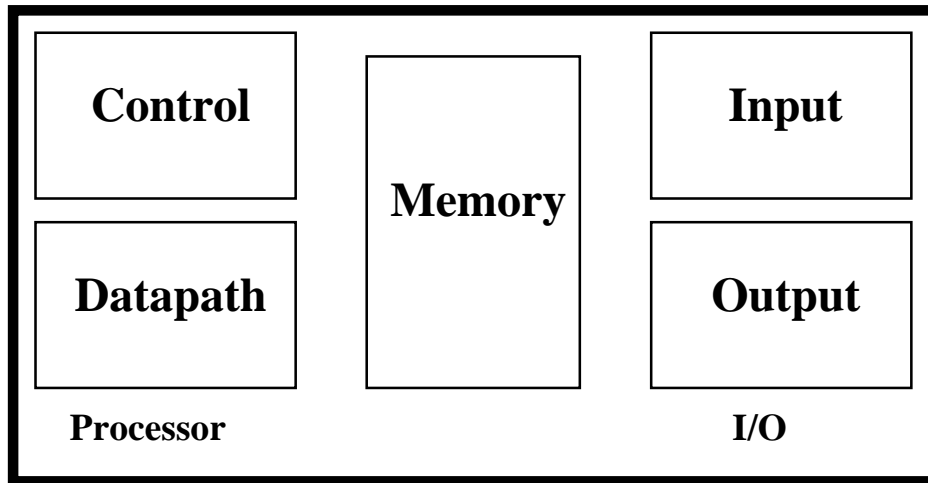
# What if we want to use only core instructions

```
#p53D.asm
# Without using pseudo instructions
#such as la "load address", li "load immediate" etc
main:
    .data    0x10010000    #starting address of first string
    .asciiz  "\nType the first number: " #msg1
    .data    0x10010100    #starting address of next
    .asciiz  "\nType the second number: " #msg2
    .data    0x10011000    #starting address of the third
    .asciiz  "\nThe answer is: "        #msg3
    .text
ori $v0, $zero, 4        #msg1
lui $a0, 0x1001          # Upper part of msg1 addr (ie a0=10010000)
ori $a0, $a0, 0          # now a0 has 1 word addr 10010000
syscall
ori $v0, $zero, 5        #read input (b)
syscall                  #now type-in is in v0
or $s1, $zero,$v0
ori $v0, $zero, 4        #msg2
lui $a0, 0x1001          #a0=10010000
ori $a0, $a0, 0x0100     #a0=10010100
syscall
ori $v0, $zero, 5        #read input (c)
syscall                  #now type-in is in v0
or $s2, $zero,$v0
add $s0, $s1, $s2        #add (b) and (c) $t for tempr
ori $v0, $zero, 4        #msg3
lui $a0, 0x1001          #a0=10010000
ori $a0, $a0, 0x1000     #a0=10011000
syscall
ori $v0, $zero,1         #request for print
or $a0, $zero, $s0      #result
syscall
```

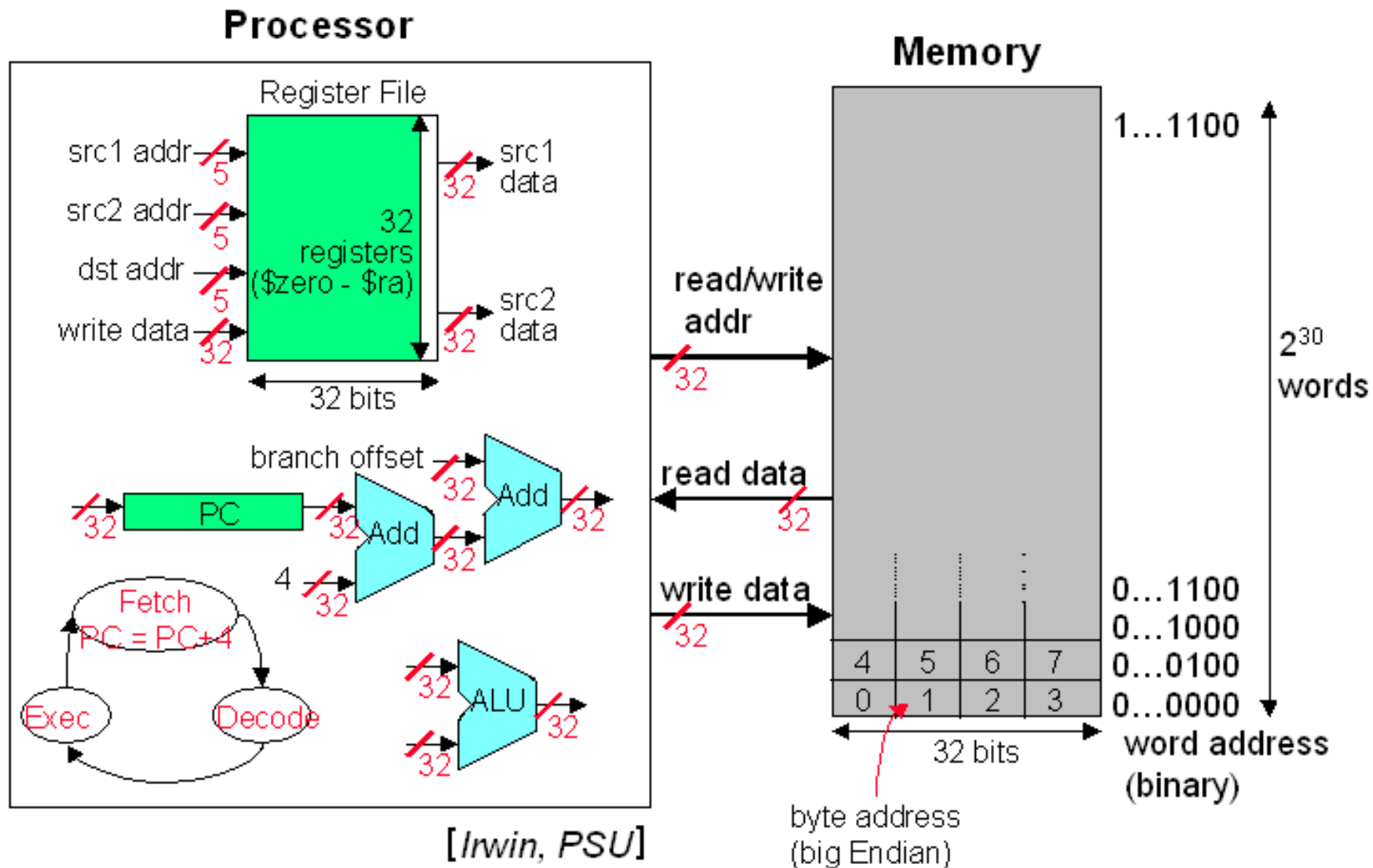


# Registers vs. Memory

- Arithmetic instructions operands must be registers,  
— only 32 registers provided
- Compiler associates variables with registers
- What about programs with lots of variables



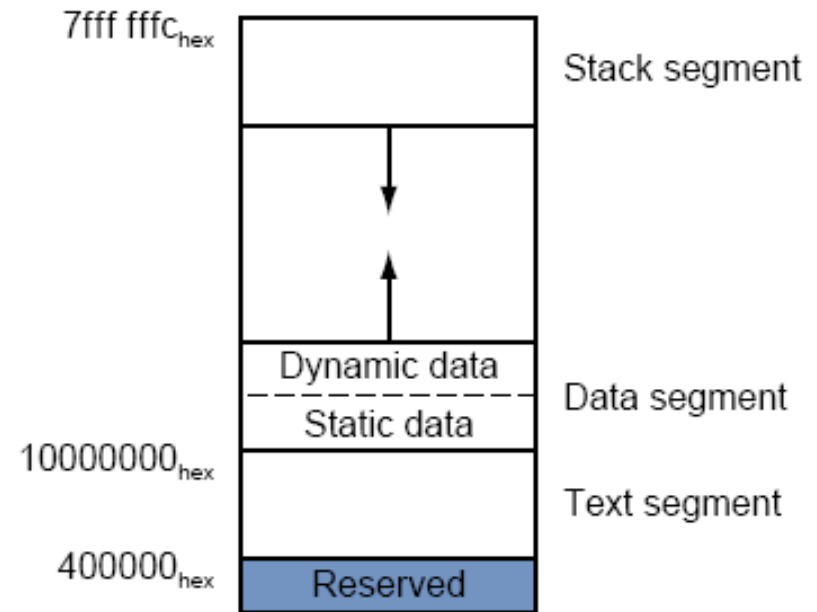
# MIPS Register & Memory



# Memory Organization

- Viewed as a large, single-dimension array, with an address.
- A memory address is an index into the array
- "Byte addressing" means that the index points to a byte of memory.

0	8 bits of data
1	8 bits of data
2	8 bits of data
3	8 bits of data
4	8 bits of data
5	8 bits of data
6	8 bits of data
...	





# Memory Organization

- Bytes are nice, but most data items use larger "words"
- For MIPS, a word is 32 bits or 4 bytes.

0	32 bits of data
4	32 bits of data
8	32 bits of data
12	32 bits of data

Registers hold 32 bits of data

...

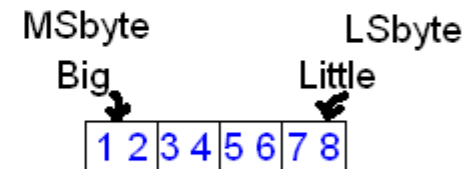
- $2^{32}$  bytes with byte addresses from 0 to  $(2^{32}-1)$
- $2^{30}$  words with byte addresses 0, 4, 8, ..  $(2^{32}-4)$
- Words are aligned  
i.e., what are the least 2 significant bits of a word address?

WORD  
0 0000 0000  
4 0000 0100  
8 0000 1000  
12 0000 1100  
16 0000 0000 ↑↑

# Byte Order and Two Camps

BIG-ENDIAN	ADDR (byte)	LITTLE-ENDIAN
Word	000F	
	000E	
	000D	
	000C	
Word	65 000B	00
	00 000A	00
	00 0009	00
	00 0008	65
Word	0007	
	0006	
	0005	
	0004	
Word	78 0003	12
	56 0002	34
	34 0001	56
	12 0000	78

"Store 0x00000065 to adress 0x0008"



"Store 0x12345678 to address 0x0000"

- Big-Endian Fill the MSbyte first at 0x0000
- Little-Endian Fill the LSbyte first at 0x0000

# Instructions for Memory Access

- **Load** and **store** instructions
- Example:

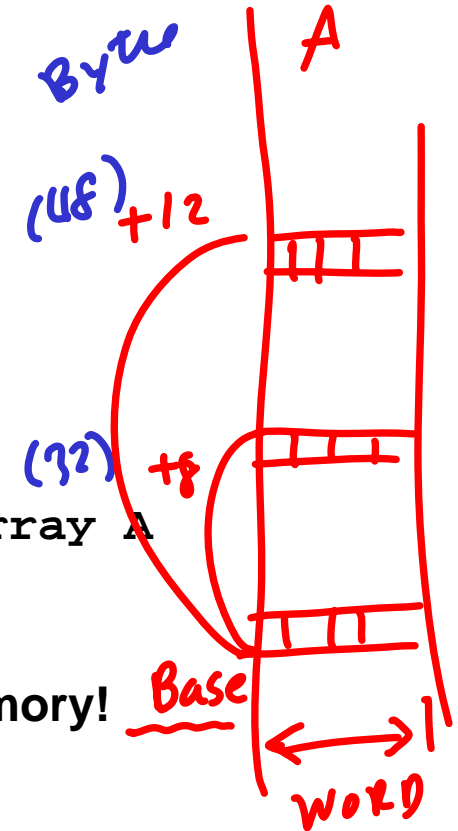
C code:            `A[12] = h + A[8];`

MIPS code:        `lw $t0, 32($s3)`  
                  `add $t0, $s2, $t0`  
                  `sw $t0, 48($s3)`

#where \$s3 holds the base address of array A

- Store word has destination last
- Remember arithmetic operands are registers, not memory!

Can't write:        `add 48($s3), $s2, 32($s3)`



# Load and Save Basics in SPIM (1/2)

```
p57.asm - Notepad
File Edit Format View Help

#p57.asm
#load and store
#the base address of array A is in $s3
#A[12]=A[1]+A[8]
main:
    .data    0x10010000          #starting address of the Array A
    .word    0x00000003, 0x00000004, 0x00000005, 0x00000006
    .word    0x00000013, 0x00000014, 0x00000015, 0x00000016
    .word    0x00000023, 0x00000024, 0x00000025, 0x00000026
    .data    0x10011000          #starting address of next
    .asciiz  "\nLoad and Store Example Program p.57\n"
    .data    0x10011100
    .asciiz  "+"
    .asciiz  "="
    .text
                                #code part
    ori $v0, $zero, 4           #msg1
    lui $a0, 0x1001             # Display the banner
    ori $a0, $a0, 0x1000
    syscall
    lui $s3, 0x1001
    ori $s3, $s3, 0
    lw $t0, 32($s3)            #t0=A[8]
#print A[8]
    ori $v0, $zero, 1           #request for print
    or $a0, $zero, $t0
    syscall
#pring + sign
    ori $v0, $zero, 4           #request for print
    lui $a0, 0x1001
    ori $a0, $a0, 0x1100
    syscall
```

# Load and Save Basics in SPIM (2/2)

```
lw $t1, 4($s3)           #t1=A[1]
#print A[1]
ori $v0, $zero, 1       #request for print
or $a0, $zero, $t1
syscall
#Operation of ADD
addu $t2, $t1, $t0
sw $t2, 48($s3)
#print = sign
ori $v0, $zero, 4       #request for print
lui $a0, 0x1001
ori $a0, $a0, 0x1102
syscall
#print the A[12]
ori $v0, $zero, 1       #request for print
lw $t3, 48($s3)         #result
or $a0, $zero, $t3
syscall
```

 Console

```
Load and Store Example Program p.57
35+4=39
```

# Check if we know this:

- MIPS
  - loading words but addressing bytes
  - arithmetic on registers only

<u>Instruction</u>	<u>Meaning</u>
<code>add \$s1, \$s2, \$s3</code>	$\$s1 = \$s2 + \$s3$
<code>sub \$s1, \$s2, \$s3</code>	$\$s1 = \$s2 - \$s3$
<code>lw \$s1, 100(\$s2)</code>	$\$s1 = \text{Memory}[\$s2+100]$
<code>sw \$s1, 100(\$s2)</code>	$\text{Memory}[\$s2+100] = \$s1$